

CSCI 210: Computer Architecture

Lecture 25: Datapath

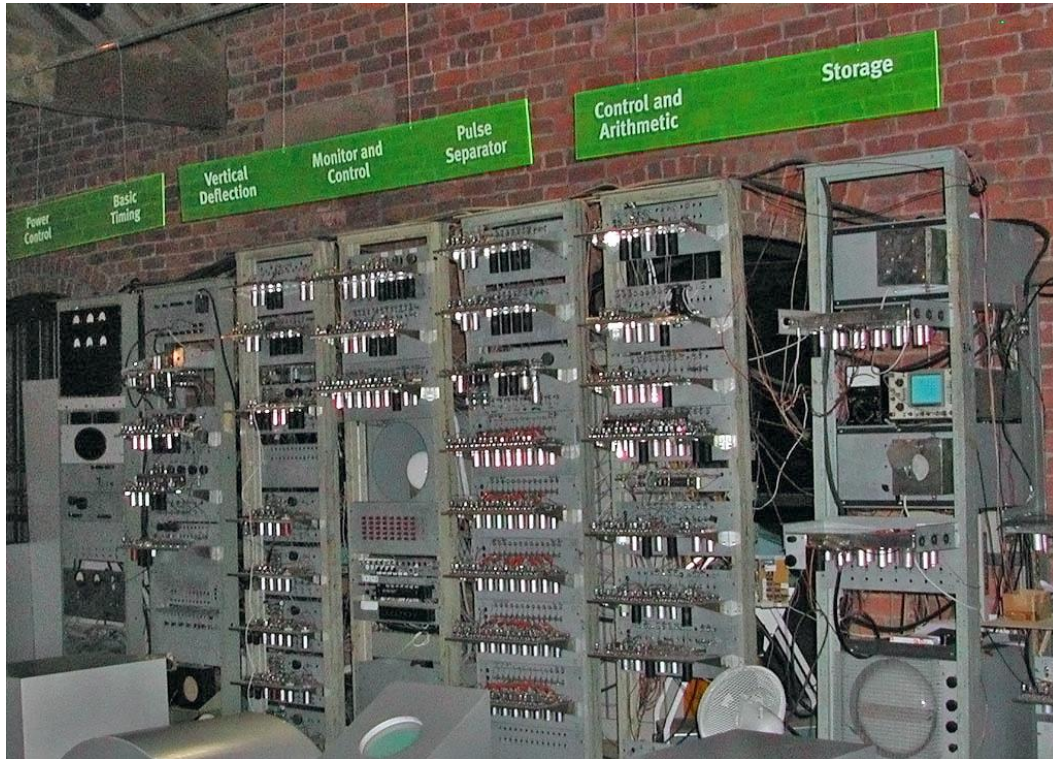
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Slides from Cynthia Taylor

Today's Class

- The datapath!

CS History: The Manchester Baby



- First stored-program computer
- Ran its first program on June 21, 1948
- Designed as a testbed for the first random-access memory
- Only arithmetic operations were addition and subtraction
- Its first program calculated the highest proper divisor of 2^{18} (262,144), by testing every integer from 2^{18} downwards
- This program was 17 instructions and took 52 minutes to run

The Processor: Datapath & Control

- We're ready to look at an implementation of MIPS simplified to contain only:
 - memory-reference instructions: `lw, sw`
 - arithmetic-logical instructions: `add, sub, and, or, slt`
 - control flow instructions: `beq`

Generic implementation

- Fetch
 - Use the program counter (PC) to supply instruction address
 - Get the instruction from memory
 - Update the program counter to the next instruction
- Decode instruction
 - Read registers
 - Read the instruction to decide how to execute
- Execute
 - Perform necessary data manipulation
 - Write to registers

To fetch an instruction, what hardware do we need?

- Fetch

- Use the program counter (PC) to supply instruction address
- Get the instruction from memory
- Update the program counter to the next instruction

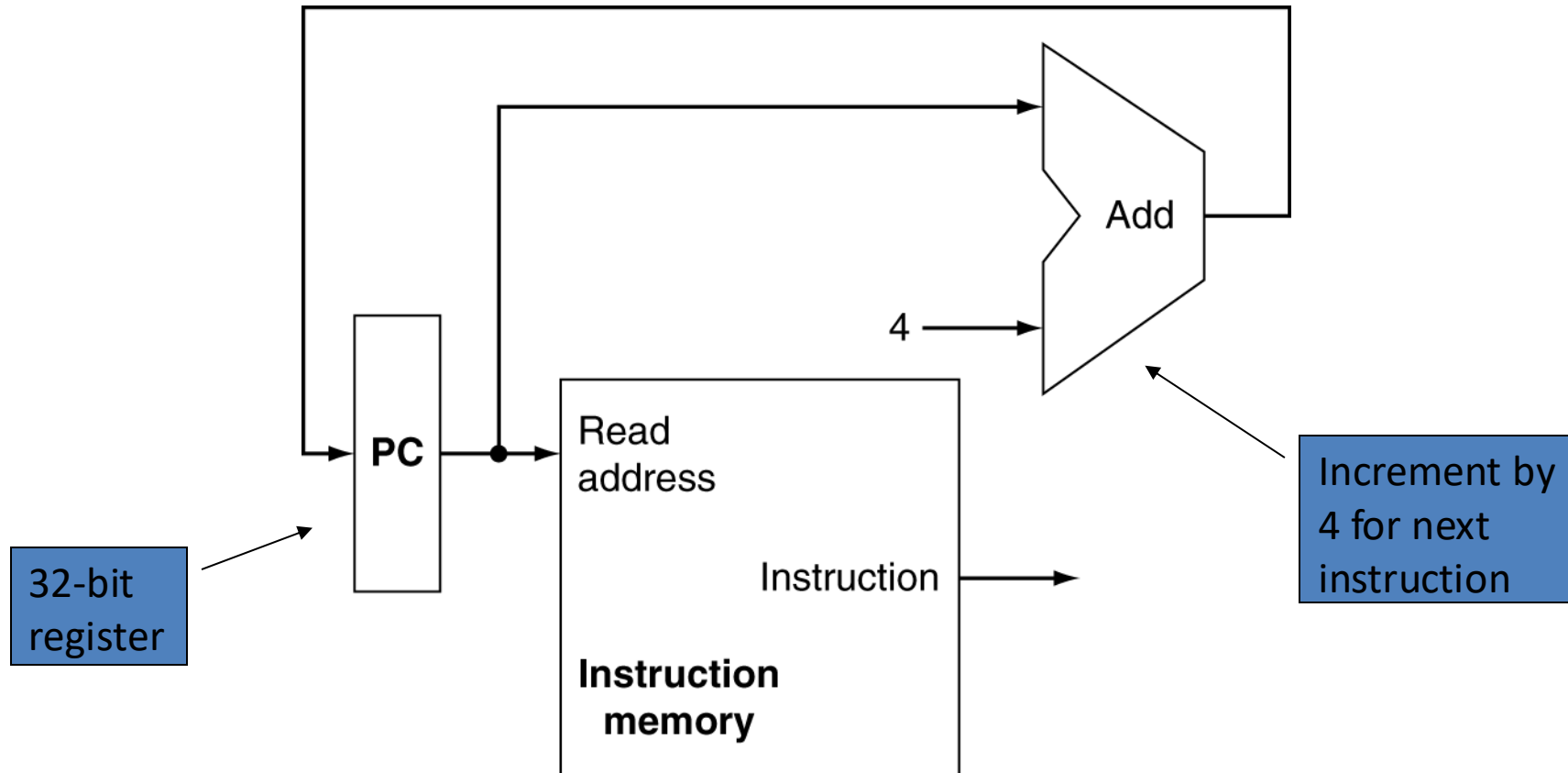
A. Register(s), Memory

B. Register(s), Adder, Memory

C. Register(s), ALU, Memory

D. More than this

Instruction Fetch



Generic implementation

- Fetch
 - Use the program counter (PC) to supply instruction address
 - Get the instruction from memory
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- Decode instruction
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Registers for instructions

- `add $t0, $t1, $t2` needs to read the values of registers `$t1` and `$t2` and write to register `$t0`
- `lw $t0, 4($t8)` needs to read one register and write one register
- `sw $t0, -8($s0)` needs to read two registers and write zero registers

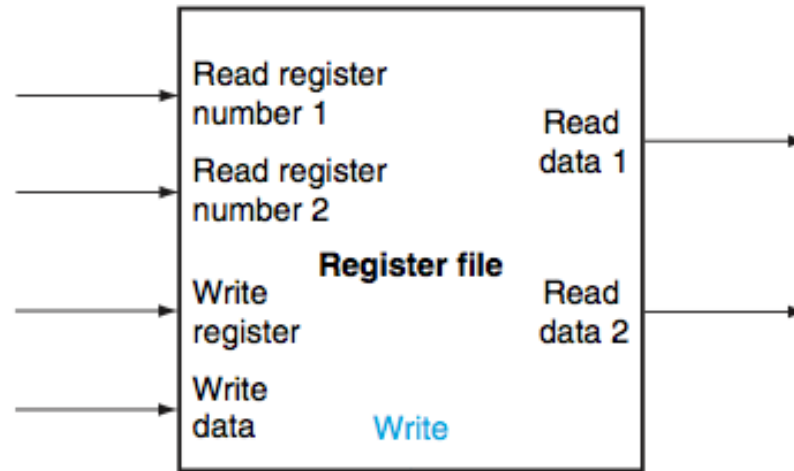
Interface for the register file

- We need the ability to read from up to 2 registers and write up to 1 register

Interface:

- Three 5-bit register select inputs (rs, rt, rd)
- Two 32-bit data outputs (data in rs and rt)
- One 32-bit data input (data to write to rd)
- One 1-bit control input (should input data be written to rd or not)

Register File



Control input that is 1 if the write data should be written to the register specified by write register

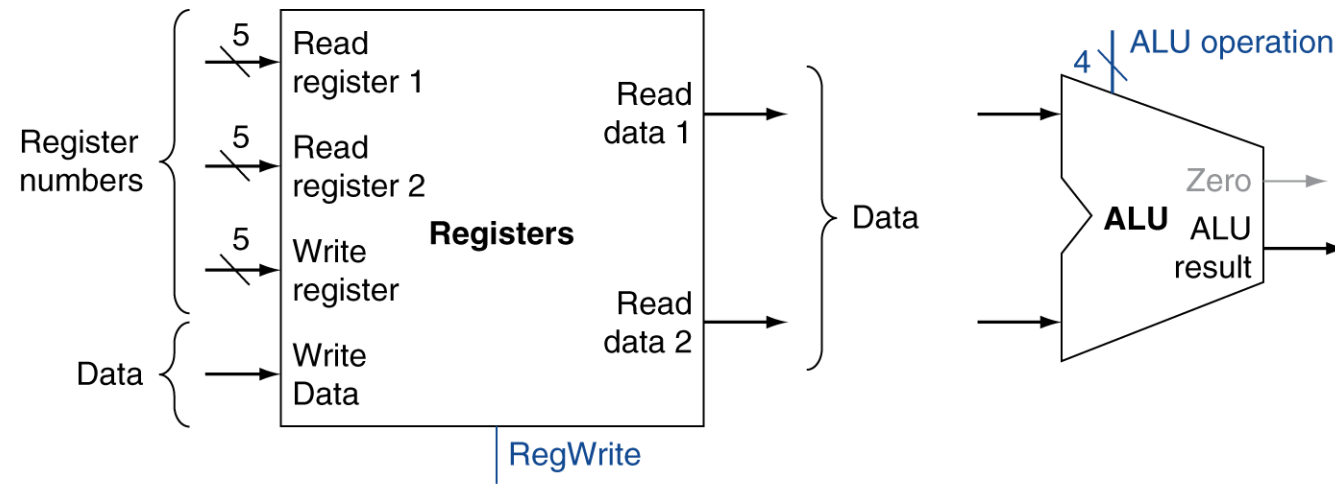


Generic implementation

- **Fetch**
 - Use the program counter (PC) to supply instruction address
 - Get the instruction from memory
 - Update the program counter to the next instruction
- **Decode instruction**
 - Read registers
 - Read the instruction to decide how to execute
- **Execute**
 - Perform necessary data manipulation
 - Write to registers

R-Format Instructions

- Read two register operands
- Perform arithmetic/logical operation
- Write register result



a. Registers

b. ALU



Data memory

- `sub $t0, $t1, $t2` does not read or write memory
- `lw $t0, 0($s0)` reads 32-bits from memory
- `sw $t0, 0($s0)` writes 32-bits to memory

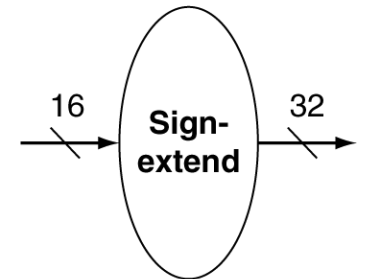
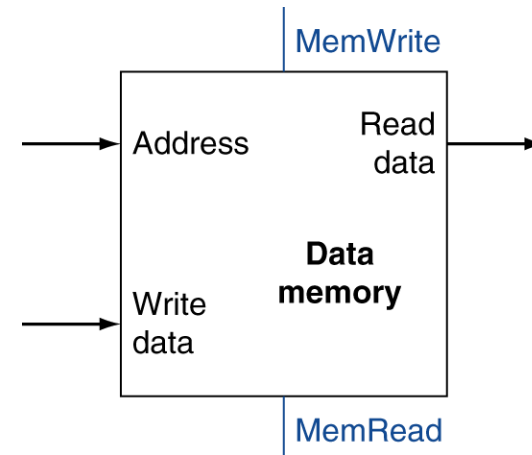
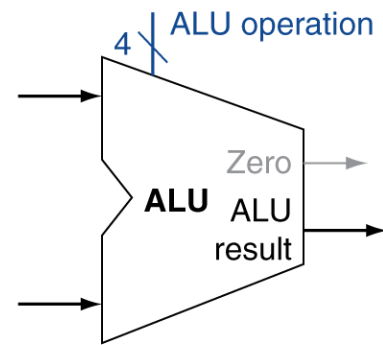
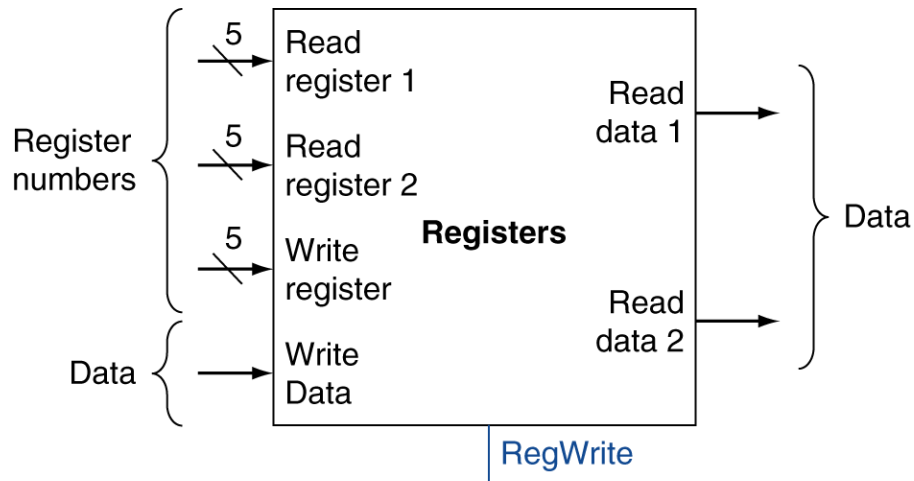
Which of these describes our interface for data memory? What do we need for lw 0(\$t3), \$t2 or sw 4(\$t5), \$t4

	Data Inputs	Data Outputs	Select inputs	Control inputs
A	One 32-bit input	One 32-bit output	One 5-bit input	2 bits
B	Zero inputs	One 32-bit output	Two 5-bit inputs	2 bits
C	One 32-bit input	One 32-bit output	One 32-bit input	2 bits
D	One 32-bit input	One 32-bit output	One 32-bit input	1 bit
E	One 32-bit input	One 32-bit output	One 5-bit input	1 bit

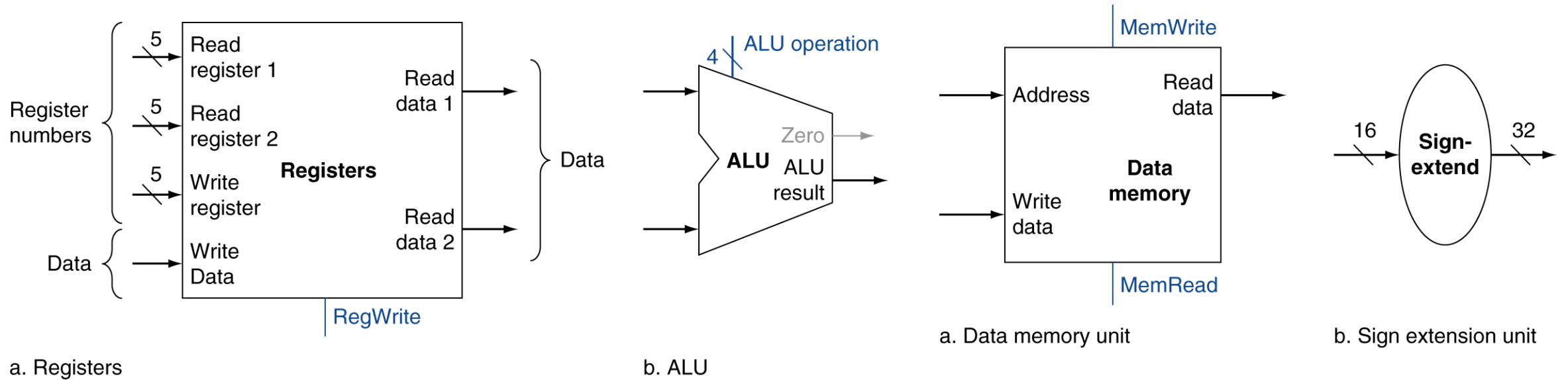
Data is what we read from/write to memory,
Select is the address we're reading/from writing to
control is what operation the data memory does (e.g., load or store)

Load/Store Instructions

- Read register operands
- Calculate address using 16-bit offset
 - Use ALU, but sign-extend offset
- Load: Read memory and update register
- Store: Write register value to memory



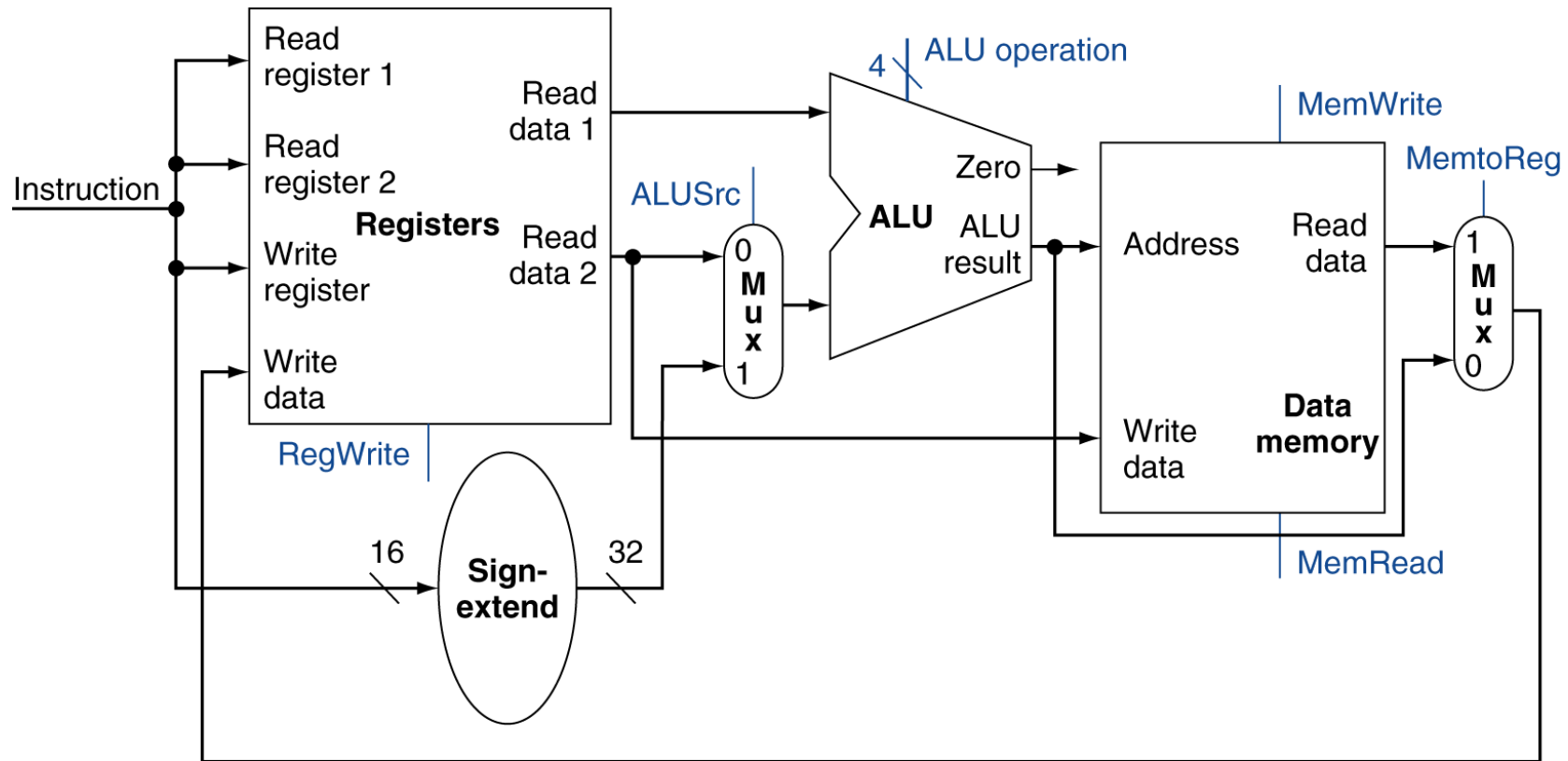
Memory



Which is true about the ALU and the register file in MIPS?

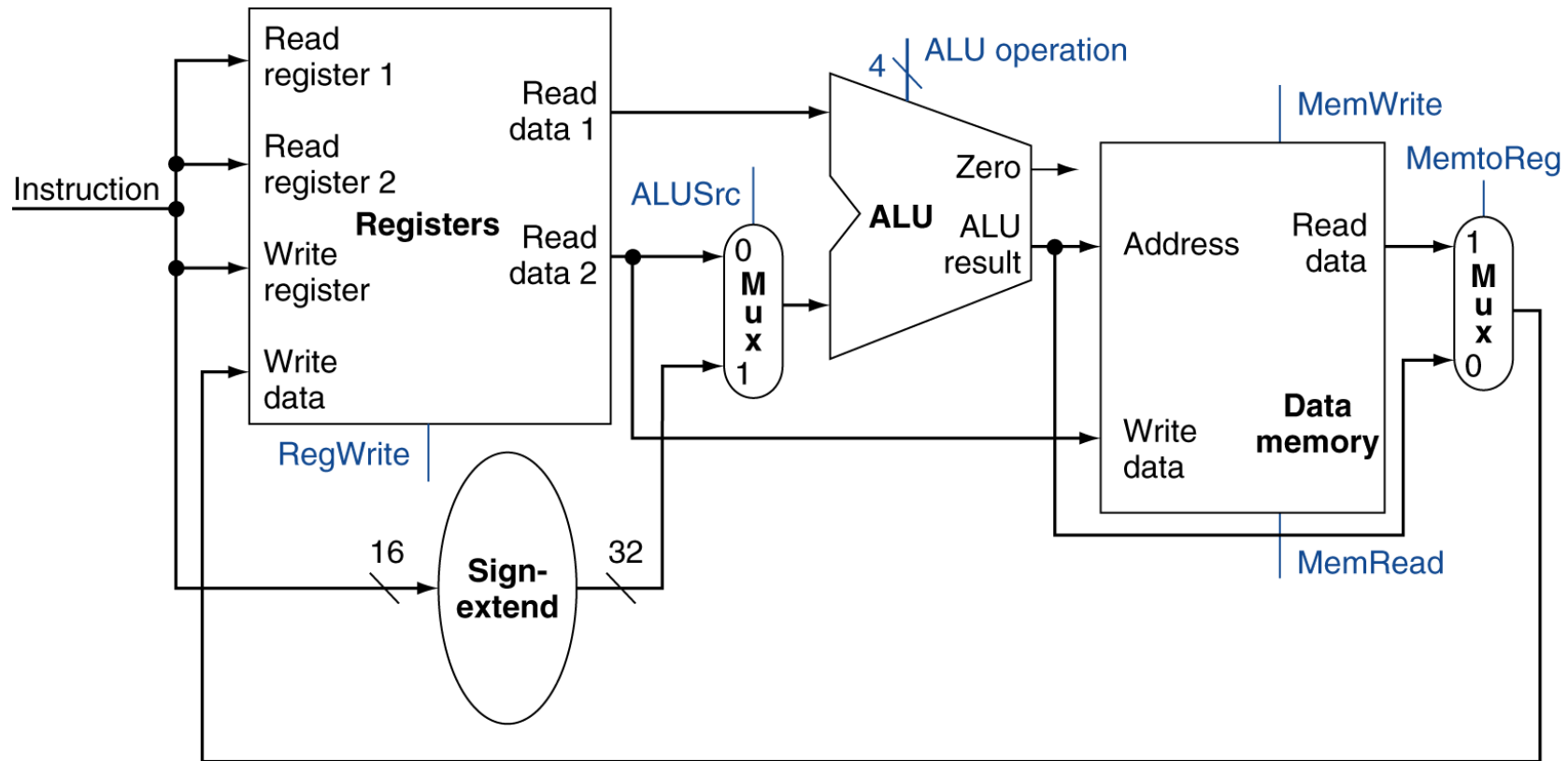
- A. The ALU *always* performs an operation before accessing the register file
- B. The ALU *sometimes* performs an operation before accessing the register file
- C. The register file is *always* accessed before performing an ALU operation
- D. The register file is *sometimes* accessed before performing an ALU operation
- E. None of the above.

R-Type/Load/Store Datapath



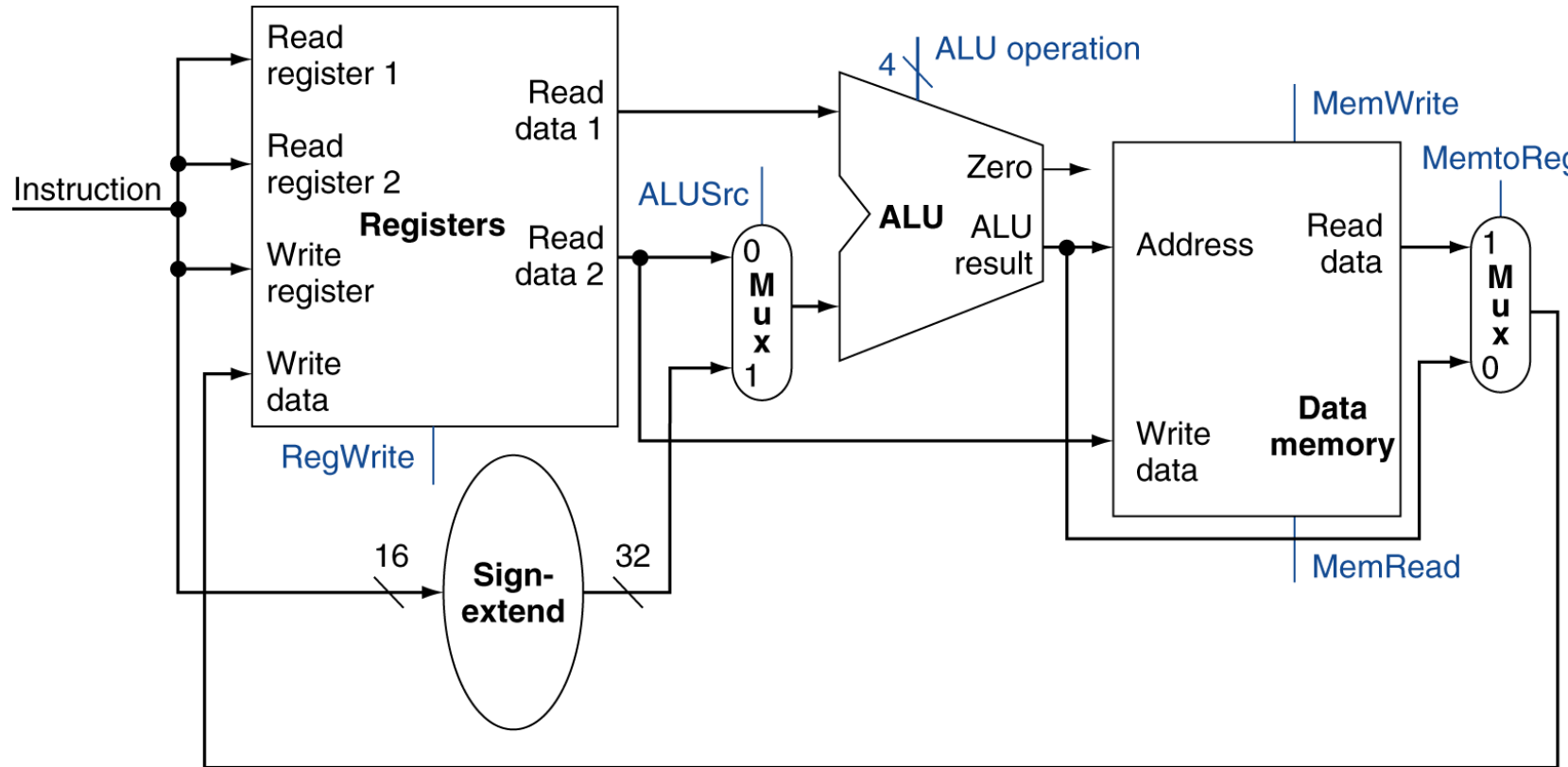
Add \$t0, \$t0, \$t1

\$t0 is register 8, \$t1 is register 9
\$t0 holds 5
\$t1 holds 6



lw \$t1, 4(\$t0)

\$t0 is register 8, \$t1 is register 9
\$t0 holds 0x07AB8110
0x07AB8114 holds 12



Branch Instructions

- Read register operands
- Compare operands
 - Use ALU, subtract and check Zero output
- Calculate target address
 - Sign-extend offset
 - Shift left 2 bits (word offset)
 - Add to PC + 4
 - Already calculated during instruction fetch

What do we need to add Conditional Branch Instructions to our datapath?

beq \$t2, \$t3, 0x4F35

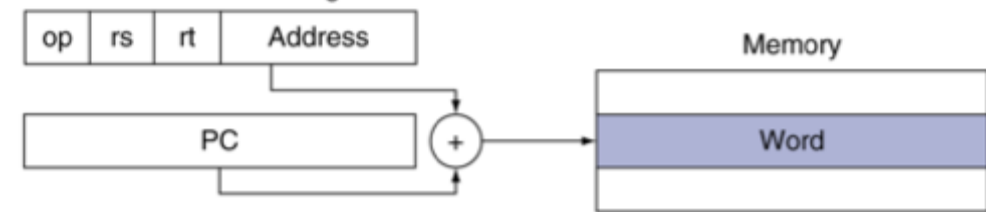
A. ALU

B. Registers and an ALU

C. Registers, ALU and Memory

D. Registers, an ALU and an Adder

4. PC-relative addressing



Read register operands

Compare operands

Use ALU, subtract and check Zero output

Calculate target address

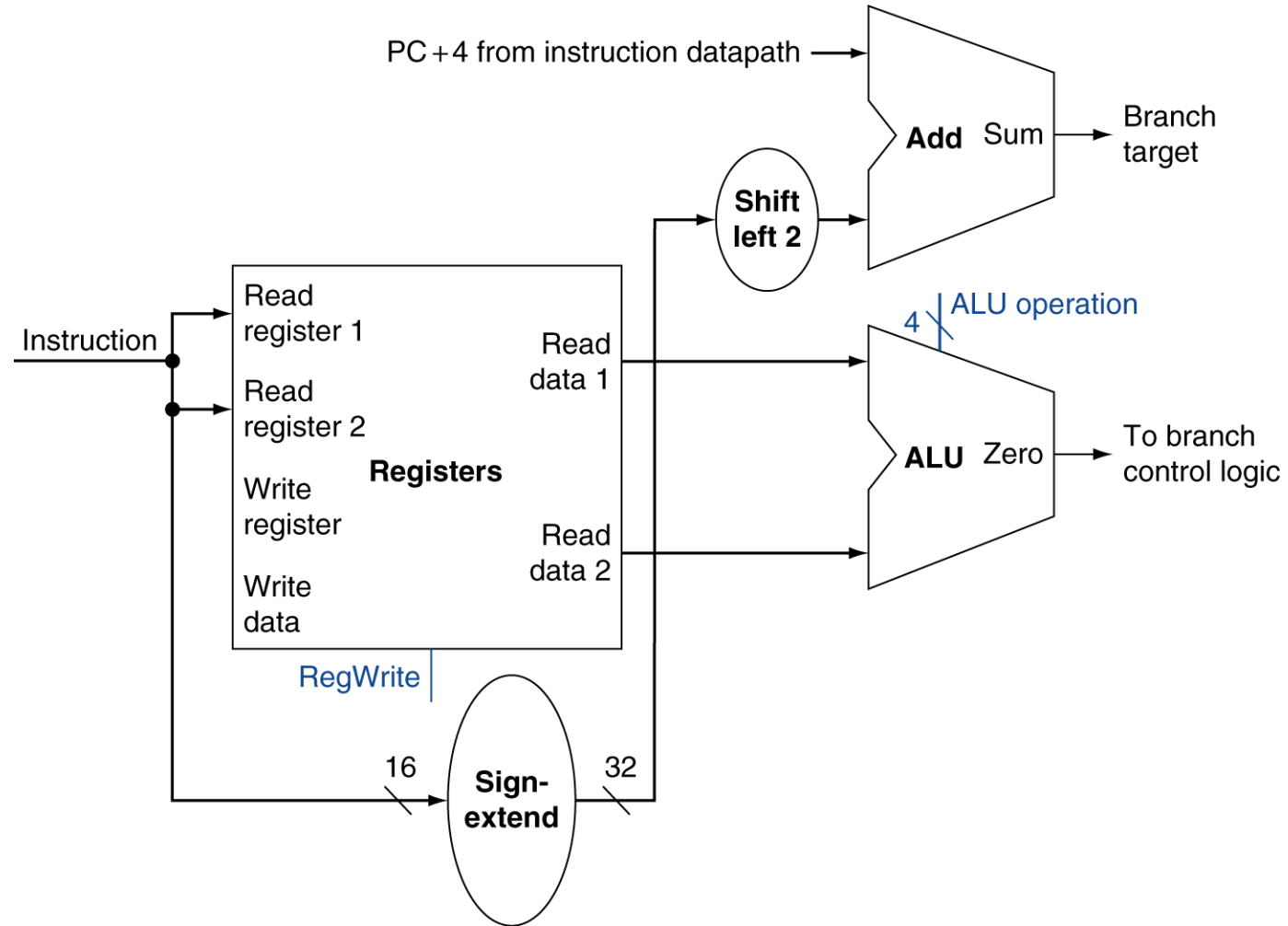
Sign-extend offset

Shift left 2 bits (word offset)

Add to PC + 4

Already calculated during instruction fetch

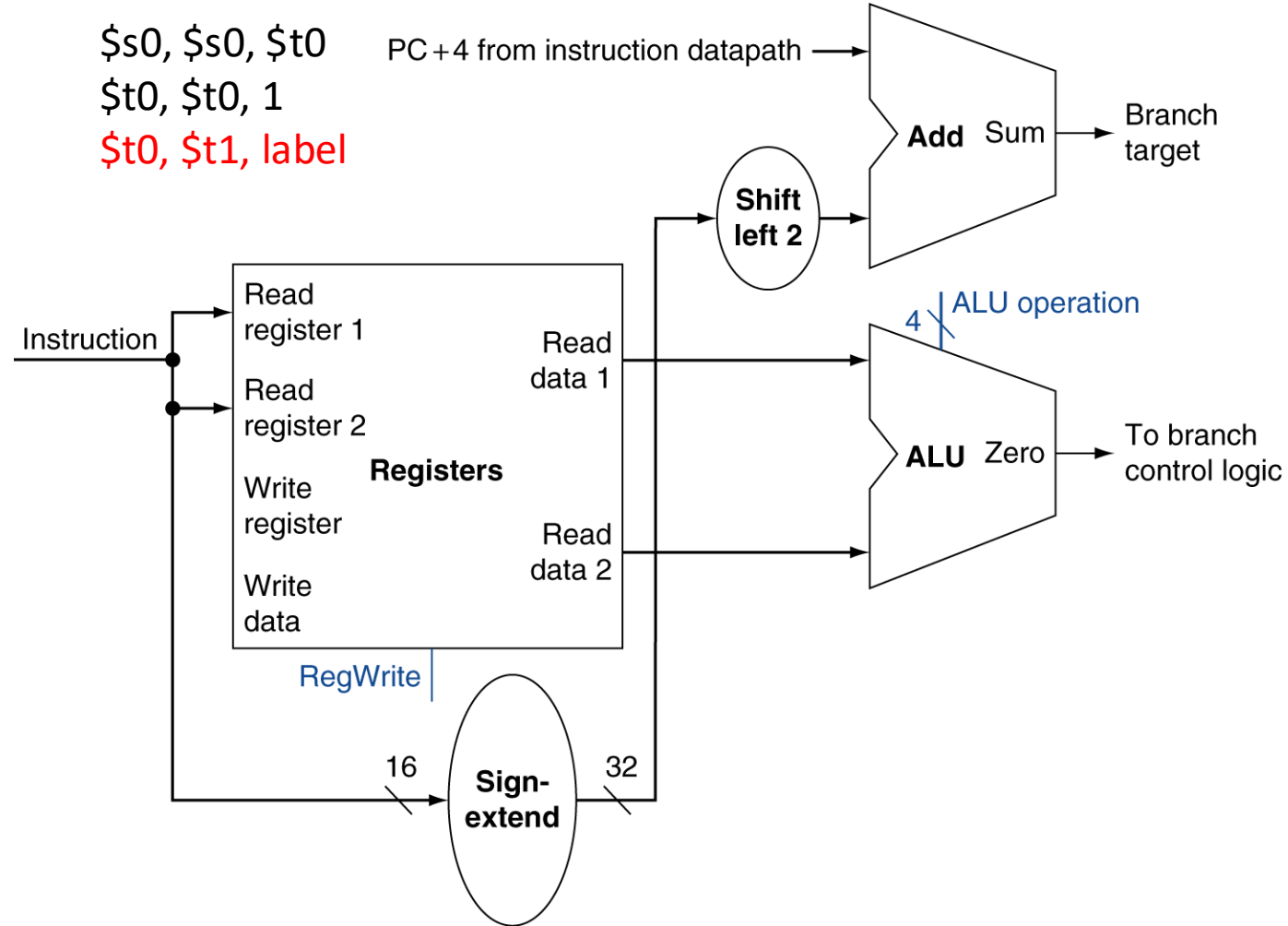
Branch Instructions



Branch Instructions

0x4045A130 label: add
 0x4045A134 addi
 0x4045A138 beq

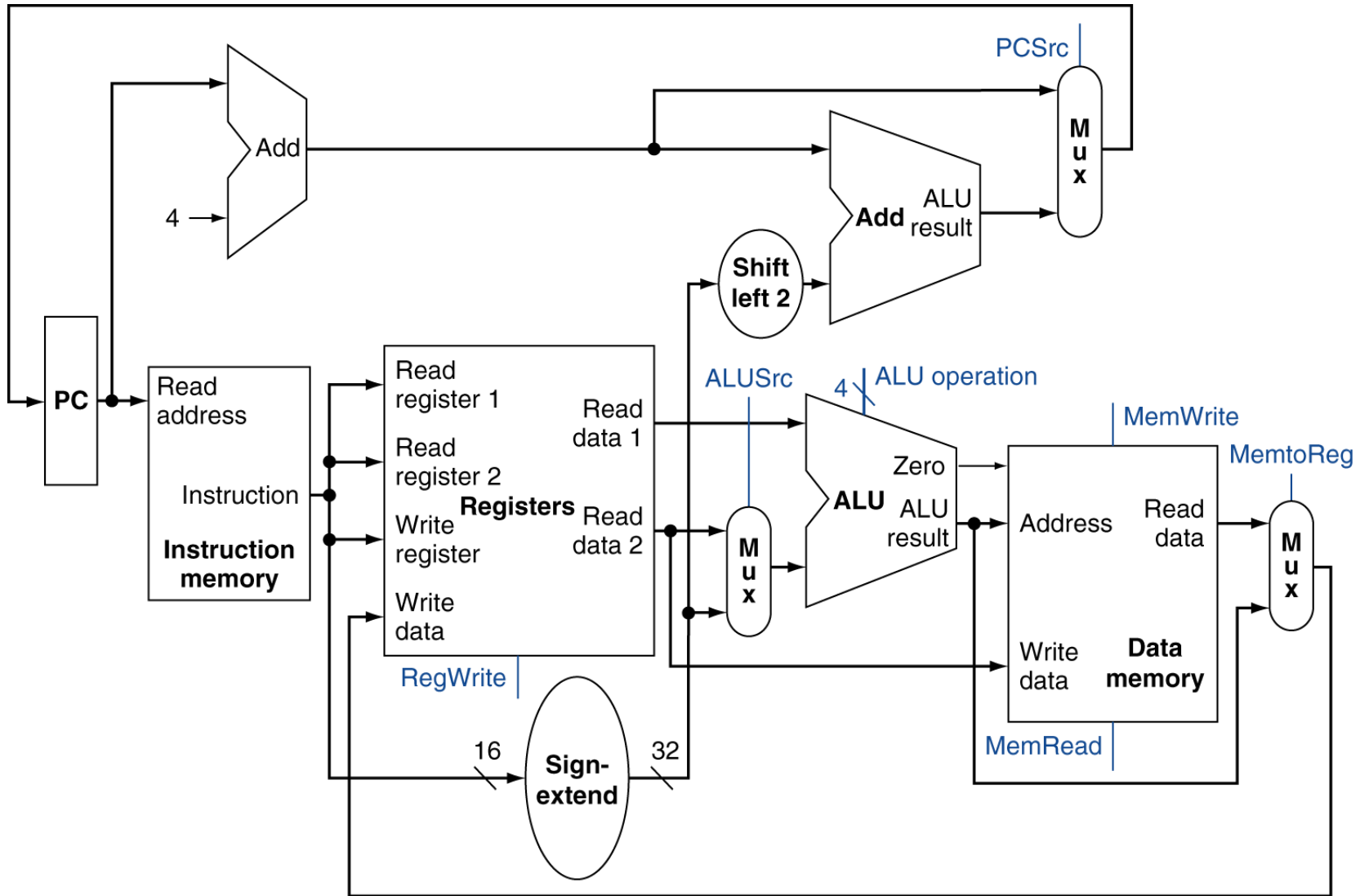
\$s0, \$s0, \$t0 PC+4 from instruction datapath
 \$t0, \$t0, 1
 \$t0, \$t1, label



op = 0x04	rs = 8	rt = 9	imm = 0xFFFFD
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\$t0 holds 5
 \$t1 holds 5

Datapath (still simplified a bit)



Reading

- Next lecture: Data path
 - Section 5.2